ANTS SYSTEMS, A STATE OF THE ART OVERVIEW: APPLICATIONS TO INDUSTRIAL PLANTS PROBLEMS

Francesco Longo, Giovanni Mirabelli, Sabrina Rondinelli
Modeling & Simulation Center, Laboratory of Enterprise Solutions (MSC-LES)
Mechanical Department, University of Calabria
Via P. Bucci, Cubo 44C, 87036 Rende, Italy
f.longo@unical.it, g.mirabelli@unical.it, srondinelli@unical.it

KEYWORDS
Ants Systems, job shop systems, Simulation

ABSTRACT
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In figure 1 ants follow a path between the nest (position in A) and the food (position in E). Suppose that, at certain point in time, an obstacle appears between points A and E. Then, ants have to decide the path to follow: AHE or ACE. As the Figure 2 shows that the two paths are not equal: route AHE is shorter than ACE. Initially ants choose the way to go random: on the average, middle of all ants goes through path AHE and the other half goes through path ACE. Subsequently the choice is influenced by the intensity of pheromone trails that ants smell on the path.

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and Dorigo 1995) who proposed several algorithms based on the same principles. While Ant’s Theory was spreading, many authors have focused on this research topic, developing algorithms for different fields of application.

Before getting into details of the state of the art overview, in the sequel a brief description of the paper sections is reported. Section 2 proposes a brief description of the main algorithms developed and the respective fields of application. In section 3 the attention is focused on algorithms applied to industrial plants problems (with a description of different cases). Section 4 discusses about one of the most important problems in industrial plants: job-shop scheduling problem. Finally, the last section summarizes conclusions and research activities still on going.

ANT’S THEORY ALGORITHMS

The first model based on Ant’s Theory was proposed by Dorigo, it was a new heuristic approach for the solution of combinatorial optimization problems called Ant System (Dorigo et al. 1991; Dorigo 1992; Dorigo et al. 1996). Ant System was first applied to solve traveling salesman problem. Over the years several versions of AS have been developed to improve and extend the performances of the initial algorithm and apply it to solve other optimization problems. In order to develop a common framework for all the versions of the AS, in 1999 was proposed the Ant Colony Optimization (ACO) meta-heuristic (Dorigo and Di Caro 1999a).

Today there are numerous successful implementations of the ACO meta-heuristic applied to different combinatorial optimization problems. It is possible to distinguish two different classes of problems: static combinatorial optimization problems and dynamic combinatorial optimization problems. Static problems are those in which the characteristics of the problem not change while the problem is being solved. On the other hand, dynamic problems are those in which the characteristics of the problem evolve during its resolution. In the following is reported a brief description of the main problems solve with ACO algorithms.

Traveling salesman problem

The first problem solved by an algorithm based on Ant’s Theory was the traveling salesman problem (TSP). TSP consists of finding the shortest route visiting each member of a collection of locations and returning to the starting point. At the beginning the algorithm used to solve TSP was the AS, in which the ants build optimum solutions (finding feasible tours of minimum costs) by moving on the problem graph from city to city. Considering the way to update the pheromone trail, three different versions of the AS have been defined: ant-density, ant-quantity and ant-cycle. In the first two versions ants deposit pheromone while building a solution, instead in third version they deposit pheromone only when a complete tour has been built. The MMAS algorithm (Max-Min AS, Stützle and Hoos 1997a; Stützle and Hoos 1997b) is similar to the AS but, in this version, the pheromone trail is updated offline. The way to update the pheromone trail is the same proposed by Dorigo and Di Caro (1999b): “…the arcs which were used by the best ant in the current iteration receive additional pheromone”. The pheromone trail values are restricted at the interval \( [\tau_{\text{min}}, \tau_{\text{max}}] \), and trails are initialized at maximum value \( \tau_{\text{max}} \).

Another version is the ASrank (Bullnheimer et al. 1997) in which, as in the MMAS version, the arcs that have been visited by ants that carried out the best tour are updates with an additional quantity of pheromone (this is equivalent to AS’s elitist ants pheromone updating).

The difference is that the ants are ranked by tour length they complete.

The Ant-Q (Dorigo and Gambardella 1995; Gambardella and Dorigo 1996) is an algorithm that tried to fuse the AS and the Q-learning (Watkins 1989). The idea of this algorithm is to update step-by-step pheromone trails with a quantity which is a prediction of the value of the next state. Suddenly it appears that the evaluation of this predictable value was extremely difficult, than Ant-Q was abandoned and replaced by ACS (Ant Colony System; Dorigo and Gambardella 1997a; Dorigo and Gambardella 1997b). The ACS gives approximately the same performances of Ant-Q but it is simpler to use. The two algorithms simply differ in the value used to update pheromone trails.

The ACS is a further version proposed to improve the performance of the AS, that gives good solutions in reasonable time only for small problems. The rationale of ACS and AS is the same but some differences exist:

- In ACS the global pheromone trails are updated offline: at the end of an iteration a quantity of pheromone is added to the arcs used by the ant that complete the best tour;
- Ants use a different decision rule, called pseudo-random-proportional rule, to choose the next city to visit;
- Ants update the pheromone trails only on-line, step-by-step, during the construction of a solution;
- ACS algorithm uses a data structure, called candidate list, that allows to have additional local heuristic information.

Quite similar to ACS is another algorithm: ACS-op-3 (Freisleben and Merz 1996a; Freisleben and Merz 1996b). The difference is that the ACS-op-3 activates a local search procedure (based on variant of the 3-opt local search procedure) to improve the solutions generated by ants.
Quadratic Assignment Problem

QAP is the problem to assign \( n \) resources to \( n \) locations trying to minimize total allocation cost (this cost is a function of the way to allocation chosen). This kind of problem is a generalization of TSP. As a consequence the algorithms used to solve a QAP are generalization of algorithms used to solve the TSP. Further information can be found in Maniezzo et al. (1994) that propose the AS-QAP algorithm and in Stützle and Hoos (1998) that propose the MMAS-QAP algorithm.

Job scheduling problem

The problem is the following: given \( M \) machines and a set of \( J \) jobs it is necessary to choose the order sequence of operations to be executed on these machines so that the maximum of the completion times of all operations is minimized and no two jobs are processed at the same time on the same machine. The algorithm used to solve this problem is the AS, but since exist some differences in nature of the constraints of this problem respect to the TSP, it is necessary to define a new way of building the ants’ tabu list. Currently AS-JSP algorithm (Colorni et al. 1994) is used to solve problems of maximum dimensions up to 15 machines and 15 jobs and the generate results are not satisfactory. This results suggest that this field of applications should be further investigated.

Vehicle Routing

Vehicle Routing problem is the following: consider a graph \( G=(N,A,d) \), where \( N \) is the set of nodes, \( A \) is the set of arcs and \( d_{ij} \) is a weight associated to arc \( (i,j) \) that represent the distance between node \( i \) and \( j \). Node 0 represents a deposit in which are localized vehicles (each one of capacity \( C \)), while the other nodes represent customers that should be served and ask a quantity \( D \) at time \( T \). The problem is to find the cycle of minimum cost such that every customer is visited and its demand dispatched respecting vehicle capacity constrains and assuring that every vehicle starts and ends its tour in the depot. The first algorithm proposed for this problem is AS-VRP that is an extension of the AS based on ASrank. Another algorithm for the Vehicle Routing problem is HAS-VRP (Gambardella et al. 1999). This algorithm first reformulates the initial problem by adding to the city a set of \( M-1 \) depots, where \( M \) is the numbers of vehicle. In this way the initial problem becomes a TSP with additional constraints. Then it is applied HAS-VRP algorithm, which is inspired by ACS. Specific research works in the field of vehicle routing have also been developed by De Sensi et al. (2006) and Curcio et al. (2007). The first research work presents an AS algorithm integrated in a simulation model for supply chain vehicle routing optimization. The second one compares the AS and Genetic Algorithms (GA) for vehicles routing optimization in a pharmaceutical distribution system.

Sequential Ordering

The sequential ordering problem consists in finding minimum weight Hamiltonian cycle on a directed graph with nodes and arcs weighted and trying to respect precedence constraints among nodes. In this case the proposed algorithm is HAS-SOP (Gambardella e Dorigo 1997), an extension of ACS. In fact, in HAS-SOP, unlike ACS, the set of feasible nodes is built taking in consideration the precedence constraints and, for local search, it is use a variant of 3-opt procedure.

Graph Colouring

Graph Colouring is the problem to find minimum number of colours for colouring a graph. The algorithm proposed is ANTCOL (Costa and Hertz 1997) that makes use of well-known graph colouring heuristics recursive large first (RLF) (Leighton 1979) and DSATUR (Brelaz 1979).

Shortest Common Supersequence

The shortest common supersequence (scs) is a common supersequence of minimal length given two sequences \( A=<a_1,…,a_m> \), \( B=<b_1,…,b_n> \). In the shortest common supersequence problem, the two sequences \( A \) and \( B \) are given and the objective is to find the shortest possible common supersequence of these sequences. In general, the scs is not unique. This problem is solved by an algorithm called AS-SCS (Michel and Middendorf 1998), which differs from AS classic because AS-SCS uses a look-ahead function which takes into account the influence of the choice of the next symbol to append at the next iteration.

Algorithms for dynamic combinatorial optimization problems

Table 1 lists all the algorithms presented in the previous sections also reporting authors and field of application. The algorithms for dynamic combinatorial optimization problems have focused on communications networks. In fact, this kind of problem is characterized by dynamic conditions. ACO implementations for communications networks are grouped in two classes: connection-oriented and connection-less networks. In the first case all the packets follow a common path in a preliminary phase of selection, instead in the second one all packets can follow different paths. Algorithms proposed for the two classes are reported in the next table (Table 2) with a list of references.
Table 1: List of applications of ACO algorithms to static combinatorial optimization problems (Dorigo and Di Caro 1999).

<table>
<thead>
<tr>
<th>Problem Name</th>
<th>Authors</th>
<th>Algorithm Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>Traveling Salesman</td>
<td>• Dorigo, Maniezzo &amp; Colorni, 1991</td>
<td>• AS</td>
</tr>
<tr>
<td></td>
<td>• Gambardella &amp; Dorigo, 1995</td>
<td>• Ant-Q</td>
</tr>
<tr>
<td></td>
<td>• Dorigo &amp; Gambardella, 1996</td>
<td>• ACS &amp; ACS-3-opt</td>
</tr>
<tr>
<td></td>
<td>• Stützle &amp; Hoos, 1997</td>
<td>• MMAS</td>
</tr>
<tr>
<td></td>
<td>• Bulsheimer, Hartl &amp; Strauss, 1997</td>
<td>• AS_rank</td>
</tr>
<tr>
<td>Quadratic Assignment</td>
<td>• Maniezzo, Colorni &amp; Dorigo, 1994</td>
<td>• AS-QAP</td>
</tr>
<tr>
<td></td>
<td>• Gambardella, Taillard &amp; Dorigo, 1997</td>
<td>• HAS-QAP</td>
</tr>
<tr>
<td></td>
<td>• Stützle &amp; Hoos, 1998</td>
<td>• MMAS-QAP</td>
</tr>
<tr>
<td></td>
<td>• Maniezzo &amp; Colorni, 1998</td>
<td>• AS-QAP</td>
</tr>
<tr>
<td></td>
<td>• Maniezzo, 1998</td>
<td>• ANTS-QAP</td>
</tr>
<tr>
<td>Job-shop Scheduling</td>
<td>• Colorni, Dorigo &amp; Maniezzo, 1994</td>
<td>• AS-JSP</td>
</tr>
<tr>
<td>Vehicle Routing</td>
<td>• Bullsheimer, Hartl &amp; Strauss, 1996</td>
<td>• AS-VRP</td>
</tr>
<tr>
<td></td>
<td>• Gambardella, Taillard &amp; Aguzzi, 1999</td>
<td>• HAS-VRP</td>
</tr>
<tr>
<td></td>
<td>• De Sensi et al., 2006</td>
<td>• AS</td>
</tr>
<tr>
<td></td>
<td>• Cucio et al., 2007</td>
<td>• AS &amp; GA</td>
</tr>
<tr>
<td>Sequential Ordering</td>
<td>• Gambardella &amp; Dorigo, 1997</td>
<td>• HAS-SOP</td>
</tr>
<tr>
<td>Graph Coloring</td>
<td>• Costa &amp; Hertz, 1997</td>
<td>• ANTCOL</td>
</tr>
<tr>
<td>Shortest Common Supersequence</td>
<td>• Michel &amp; Middendorf, 1998</td>
<td>• AS-SCS</td>
</tr>
</tbody>
</table>

Table 2: List of applications of ACO algorithms to dynamic combinatorial optimization problems (Dorigo and Di Carlo 1999)

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<tr>
<th>Problem Name</th>
<th>Authors</th>
<th>Algorithm Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>Connection-oriented network</td>
<td>• Schoonderwoerd, Holland, Brute &amp;</td>
<td>• ABC</td>
</tr>
<tr>
<td>routing</td>
<td>Rothkranz, 1996</td>
<td></td>
</tr>
<tr>
<td></td>
<td>• White, Pagurek &amp; Oppacher, 1998</td>
<td>• ASGA</td>
</tr>
<tr>
<td></td>
<td>• Di Carlo &amp; Dorigo, 1998</td>
<td>• AntNet-FA</td>
</tr>
<tr>
<td></td>
<td>• Bonabeau, Herma, Guérin, Snyers, Kuntz &amp; Théraulaz, 1998</td>
<td>• ABC-smartants</td>
</tr>
<tr>
<td>Connection-less network</td>
<td>• Di Carlo &amp; Dorigo, 1997</td>
<td>• AntNet &amp; AntNet-FA</td>
</tr>
<tr>
<td>routing</td>
<td>• Subramanian, Druschel &amp; Chen, 1997</td>
<td>• Regular ants</td>
</tr>
<tr>
<td></td>
<td>• Houes, Guérin, Snyers &amp; Kuntz, 1998</td>
<td>• CAF</td>
</tr>
<tr>
<td></td>
<td>• Van der Put &amp; Rothkranz, 1998</td>
<td>• ABC backward</td>
</tr>
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The next section focuses on ants’ algorithms developed for specific applications in industrial field.

**ANT’S THEORY APPLIED TO PRODUCTION SCHEDULINGS IN INDUSTRIAL PLANTS**

The problem of production scheduling within industrial plants consists of scheduling jobs (i.e. shop orders) on the available machines over the time. The production scheduling problem allows multilevel classification:

1) Depending on the number of machines:
   - Scheduling on single machine: in this case it is necessary to find the sequence of jobs that minimize use time of a single machine (in literature known as “problems of pure scheduling);
   - Scheduling on multi-machines: in this case jobs should be processed on more machines.
2) The scheduling on multi-machines can be also divided in:
   - Scheduling on homogeneous machines: all machines execute the same operation and jobs can be processed indifferently on each machine. The problem is twofold: assignment of each job to one of the alternative machines (first level: assignment problem); jobs scheduling on each machine (second level: scheduling problem), with the aim to optimising an objective function.
   - Scheduling on heterogeneous machines: each machine executes a particular operation.
3) For jobs scheduling the following case have to be considered:
   - Open shop: jobs do not have a fixed machines visiting order;
   - Flow shop: jobs have to be processed through a fixed sequence of machines and the processing operation order of each job is the same.
   - Job shop: jobs have to be processed through a sequence of machines, but the processing operation order as well as the sequence of machines of each job is different from other.

The following figure 3 represents schematically the previous scheduling problems classification within industrial plants. In the following we focus the attention first on single machine problem and then on the last three problems (open, flow and job shop) and, in particular, we present the ACO algorithms proposed to solve these problems.

**Figure 3: Classification of production scheduling problems.**

### Single Machine

In this case jobs utilize the same resource (machine). In a real world system (i.e. an industrial plant) several
elements make this problem very complex: process time, due date, and above all setup time. In literature, we find that this problem has been classified as NP-hard and its description is as follows: there are n jobs to be processed on the same machine; the machine can process only one job at a time. Each job is characterized by: processing time, due date, weight in terms of cost for time unit and setup time. To solve this problem an ACO heuristic has been presented by Liao & Juan (Liao & Juan 2005). The ACO algorithm proposed by Liao & Juan is basically the ACS version but it introduces the minimum pheromone value from “max-min-ant-system” (MMAS). The algorithm has two additional distinctive features: it introduces a new parameter for the initial pheromone trail and adjusts the timing of applying local search.

Another ACO algorithm for the same problem was proposed by Gravel (Gravel et al. 2001). In this case the algorithm is developed on a real industrial case study: a scheduling problem in an aicam aluminum foundry. In this algorithm the ACO meta-heuristic allows to perform a multi-objective optimization: minimization of unused capacity, minimization of the total tardiness of the set of orders, minimization of the total number of draining required when changing materials, minimization of transportation function penalty that computes the total unused vehicle capacity. To adapt the algorithm the authors implement a distance matrix D that contains information derived from each of the objectives and, on the other hand, a new rule for updating the pheromone trail at the end of a cycle based on one of the objectives.

Open Shop

The Open Shop problem can be defined as follows: a finite set of jobs, O, that have to be processed on a set of machines, M. Jobs do not have a fixed order to visit machines and each machine can process one job at a time. The objective is to minimize the makespan (the maximum of jobs completion time).

The algorithm proposed to solve this problem, based on Ant’s Theory, is Beam-ACO (Blum, 2005). This algorithm is an union between Ant Colony Optimization and Beam Search (Ow PS 1988), a classical tree search method that is a derivation of branch and bound algorithm. Beam-ACO works as follows: for each iteration, n ants perform a probabilistic beam search and each ant constructs up to k solutions. Then, depending on a measure, called convergence factor, and a boolean control variable, an update of the pheromone values is performed. After, the convergence factor is recomputed and it is decided if the algorithm has to be restarted. The algorithm stops when the termination conditions are satisfied.

Flow Shop

The simplest case of Flow Shop is the 2-machine flowshop scheduling problem with n jobs to schedule on 2 machines. Every jobs have a specific processing time on the two machines (i.e. job i has processing time \( a_i \) on first machine and \( b_i \) on second machine). \( C_i \) is the completion time of job i on machine 2, where \( i=1...n \).

For these types of problems two criterions can be defined: (i) the makespan criterion, called \( C_{max} \), defined as the maximum completion time of jobs on machine 2; (ii) the total completion time criterion, called \( \Sigma C_i \), defined as the sum of completion time of jobs on machine 2. It is supposed that the total completion time criterion has to be minimized subject to the condition that the makespan computed is minimum. It is supposed to minimize the makespan (maximum completion time) as the primary objective and minimize total flow-time (sum of the completion times of all jobs) as the secondary objective. In practice this problem is a multiple criteria scheduling problem and it is a strongly NP-hard.

The 2-machine flowshop scheduling problem has been extensively studied in literature and both exact and heuristic algorithms have been proposed. An ACO algorithm was also proposed to solve this problem. The ACO heuristic proposed for this problem is known as SACO (T’kindt et al. 2001). In SACO each ant builds a feasible schedule by using a constructive procedure. This procedure uses a particular structure called “pheromone matrix” that, based on memory shared by ants, contains the information about the ways to build a good feasible schedule. Every element that composes the pheromone matrix represents a probability (i.e. \( \tau_{ij} \) is the probability of having job i at position j in a good schedule for the \( \Sigma C_i \) criterion).

Starting from position 1 to position n, the most suitable job for position j is chosen according to either the intensification mode (the ant chooses the job with the highest value of \( \tau_{ij} \)) or the diversification mode (the ant uses a wheel process to select the most suitable job). When an ant has built a complete schedule, a local search is applied, that among various schedules keeps the schedule that has an optimal value of the makespan and the lowest value for the total completion time. Several computational experiments have suggested that SACO yields better results than existing heuristics.

Another variant of 2-machine flowshop scheduling problem involve the machine setup time. In this case the objective function is to find a schedule of jobs such that the total completion time, the sum of completion times of all jobs, is minimized. To solve this problem was proposed a particular variant of ACO (Shyu et al. 2004). First of all the original problem is transformed into a graph-based-model in which every edge has an associated weight \( w \). Also, unlike the strategy used by Dorigo in classical ACO, where every edge has a constant quantity of initial pheromone, this method initializes the pheromone on edges on the basis of results obtained by a greedy heuristic algorithm. To obtain a better initialization of pheromone intensity, the greedy heuristic is applied \( n \) times by letting every node be the starting node. Obtained the initial solution the
algorithm starts choosing the next node to visit and updating pheromone trail. The algorithm ends after a given number of iterative cycles. The flow shop problem becomes more complicate in case of \( n \) machines. Rajendran & Ziegler (2002) propose two algorithms to solve this problem. The first algorithm, called M-MMAS, extends the ant-colony algorithm called MMAS, by incorporating the summation rule developed by Merkle and Middendorf (2000) for the single-machine total weighted-tardiness problem, and modifying the procedure for job selection to be appended to the partial ant-sequence. In addition, the M-MMAS makes use of a new local search technique, called job-index-based local search. The second algorithm developed by Rajendran & Ziegler (2002) is called PACO. The first sequence of the ant-colony algorithm is obtained in a way similar to M-MMAS with the objective function value set to best value. There is a specific procedure to select an unscheduled job \( j \) for position \( k \). The resulting sequence is subjected to the job index-based local search scheme three times to improve the solution. Moreover the algorithm has a procedure to update pheromone trail based, not only on the resultant sequence obtained after the three-time application of the job index-based local search procedure, but also on the relative distance between a given position and the position of job \( i \) in the resultant sequence. PACO algorithm generates 40 ant-sequences and among them it is obtained the best heuristic sequence.

Another important contribution is the application of ACS to solve flow shop problem, in which \( n \) jobs should be processed on \( m \) machines, with the objective to find a permutation of jobs that minimizes the makespan. The problem can be represented in ACS by a disjunctive graph, where the set \( O \) contains all operations of each job on machine, \( C \) contains the precedence relationships between the processing operations of a single job, and \( D \) represents the machine constraints of operations belonging to different jobs. Also, there are two additional nodes: a nest node \( N \) and a source food node \( F \). After the construction of the graph, since the jobs have all the same ordering sequence, it is simple to find the first sequence. The next node an ant chooses to move to is calculated by applying a specific state transition rule. The chosen node is then added to the tabu list and the process is iterated. At the end, the node permutation given by the tabu list can determine the job sequence.

Another algorithm to solve the flow shop scheduling problem was proposed by Gajpal et al. (2005), called NACO. This time the aim is to minimize the variance of jobs completion times (called the completion-time-variance problem, CTV). In the NACO algorithm the initial solution is obtained by NEH (Nawaz, Enscore and Ham) heuristic considering the jobs CTV minimization. To improve the initial sequence is used a random-job-insertion local search procedure and for every ant-sequence generated and for final solution is executed the local search procedure three times. As in every ACO algorithm, there is a step to update the pheromone trail or trail intensity.

**Job Shop**

Among different scheduling problems previously presented the job shop is the most difficult and problematic to handle. Nowadays most of the industries, developing FMS (Flexible Manufacturing Systems), have to deal with job shop scheduling problem. As a consequence numerous algorithms and procedures have been developed over the time to handle this problem. In the last years also algorithms based on ACO meta-heuristic have been proposed. The first algorithm handles the problem of FJP (flexible job shop scheduling) which extends the classic problem of job-shop scheduling. In fact, in a very large number of industries all available machines are divided into groups of identical machine tools. This version of job shop scheduling problem is known as FJS with parallel machines. A scheduling algorithm was introduced by Rossi and Dini (2006) and represents the FJS problem with separable transportation and sequence dependent setup times in a disjunctive graph. The aim of the algorithm is to minimize the makespan. The algorithm proposed is ACS with a number of innovative skills: the first is an approach based on the disjunctive graph model and a LS (List Scheduler) algorithm (that drives an ant to visit the digraph model). The second feature is the pheromone trail structure that is based on the before mentioned disjunctive graph model. Finally, two original components, the routing-precedence-based visibility function and the method to approximate non-delay schedules are introduced to improve the performance.

For FJS problem another important procedure have been obtained combining GA (genetic algorithm) with ACO (Rossi & Boschi 2008) as a mechanism of collaborative search. This procedure presents new skills in the methods of selecting the GA subpopulations and updating pheromone trails. This algorithm has been implemented in a real system to test its performances; the results show that the algorithm is a powerful mechanism for enhancing the searching capability of both genetic algorithms and ant colony systems.

**ACO AND MODELLING AND SIMULATION**

This paragraph discusses about simulation use in job shop scheduling problem (in real industrial environments) and in particular about the implementation of simulation models supported by ACO meta-heuristic. The job shop scheduling problem within real industrial environments is usually characterized by a number of restrictive assumptions (i.e. complete resources utilization, no preemption, a job cannot simultaneously occupy more than one machine at a time, etc.). Restrictive assumptions generate ideal scenarios, useful
to gain confidence and knowledge about the real system, but they do not give the possibility to transfer and apply final results to the real system. Usually restrictive assumptions can be deleted using Modeling & Simulation to recreate systems complexity.

In the following we report an example of simulation model to solve the flexible job shop scheduling problem. This work, proposed by Li-Ning Xing et al. (2008), considers the following restrictive assumptions for the job shop scheduling problem:

- Jobs are independent from each other;
- Machines are independent from each other;
- Setting up time of machines is negligible;
- Move time between operations is negligible;
- At a given time, a machine can execute at most one operation;
- No more than one operation of the same job can be executed at a time;
- There are no precedence constraints among the operations of different jobs.

The aim of a simulation model is to assign the operations to machines and to schedule the operations on each machine according to certain rules (i.e. scheduling rules) or algorithms (i.e. genetic algorithms, ACO algorithms).

In the past, the authors have already introduced Modeling & Simulation (M&S) for supporting production scheduling in real industrial plants (in particular job shop systems). Longo et al. (2003) presents a simulation model for shop orders scheduling in special glass manufacturing; Longo et al. (2006-a) use M&S combined with scheduling rules for production planning within a manufacturing plant devoted to produce high pressure hydraulic hoses. Longo et al. (2006-b) face the scheduling problems (and plant layout optimization) in a wood frames manufacturing process. Finally Longo et al. (2008) combine M&S and genetic algorithms for shop orders scheduling in a real manufacturing plant.

AN APPLICATION EXAMPLE

This section presents the initial results of a research work (developed by authors) that aims at using Modeling & Simulation in a job shop scheduling problem supported by ACO meta-heuristic.

The job shop system manufactures high pressure hydraulic hoses (Longo et al., 2006-a). Each product is made up by a high pressure hose, two adapters and two hydraulic fittings. The production process is made up by 8 different operations:

- Materials preparation
- Fittings stamp
- Hoses cutting
- Hose skinning
- Assembly
- Junction
- Hydraulic hoses Testing
- Final controls and packaging

Two types of Shop Orders (S.Os) can enter the system: normal S.Os and priority S.Os. Usually normal S.Os are scheduled on a 2-weeks time window (each new S.O. enters in the last position of the 2-weeks queue). On the contrary, a priority S.O. can enter the 2-weeks queue in any position at any time (it depends on the priority level of the S.O.). In other words the system allows the passing between jobs. Each S.O. has a finite number $m$ of operations, one on each machine and it is allowed to work twice a job on the same machine. All the S.Os entered into the system must be necessarily completed.

Machines could not be available during the scheduling period because of the failures. Failures have been modeled by using a negative exponential distribution for both the Mean Time to Failure (MTTF, expressing the time between two consecutive machine failures) and the Mean Time to Repair (MTTR, expressing the time required for repairing the machine). Finally process and set-up’s times are considered as stochastic variables each one with a specific distribution form based on historical data. According to these hypotheses it follows that the case analysed belong to the dynamic-stochastic scheduling problem because new S.Os arrive during the scheduling horizon and most of the numerical quantities are stochastic.

The simulation model of the job shop system was developed by using the discrete event simulation software eM-Plant. The figure 4 shows the main frame of the simulation model.

![Fig. 4 – Simulation model main frame](image-url)

Further implementations have been carried out to introduce an Ant System algorithm as support tool for S.Os scheduling. The implementation of the algorithm as well as the interface with the simulation model was created through the programming of specific subroutines, written using the simulation language Simple++.

Some initial optimizations have been carried out trying to minimize the S.Os flow time and lateness and maximize the fill rate (the last one intended as percentage of S.Os that meet the due date). The flow time reduction was about 10.5% lower than the best result obtained by using the Short Production Time (SPT) scheduling rule. Analogously the lateness...
CONCLUSIONS

The paper presents a state of art about Ants System with particular attention to industrial plants problems. The state of art shows that the main fields of application for ACO algorithms are the routing problems, but in the last years, different applications have been developed to solve industrial plants problems and, in particular, production scheduling problems (job shop, flow shop, open shop). The final part of the paper focuses on the integration of Modeling & Simulation and Ant Theory for supporting real time production scheduling in job shop systems. The authors present some initial results on the development of a simulation model of a real industrial plant (job shop system) supported by an AS algorithm. Further researches are still on going for simulation model accreditation within the real system in order to use it a support tool for real-time production scheduling.

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FRANCESCO LONGO took the degree in Mechanical Engineering from University of Calabria (2002) and the PhD in Industrial Engineering (2005). He is currently scientific responsible of the MSC-LES (Modeling & Simulation Center – Laboratory of Enterprise Solutions). He is also Associate Editors of the international journal Simulation: Transaction of SCS. He supports the organization of the most important conferences on M&S (SCSC, EMSS, MAS, etc.). His research interests include modeling & simulation of manufacturing systems and supply chain management, vulnerability and resilience, DOE, ANOVA.

GIOVANNI MIRABELLI was born in Rende (Italy), on January the 24th, 1963. He took the degree in Industrial Technology Engineering from University of Calabria. He is currently researcher at the Mechanical Department (Industrial Engineering Section) of University of Calabria. His research interests include Modeling & Simulation for workstations effective design, work measurement and human reliability.

SABRINA RONDINELLI was born in Tropea (Italy) in May the 9th, 1984. She took her degree in Management Engineering, summa cum Laude, in September 2008 from the University of Calabria. She is currently PhD student at the Mechanical Department of University of Calabria. Her research activities concern the development of simulation models, based on general purpose programming languages and artificial intelligence techniques, for production and Supply Chain planning and management.